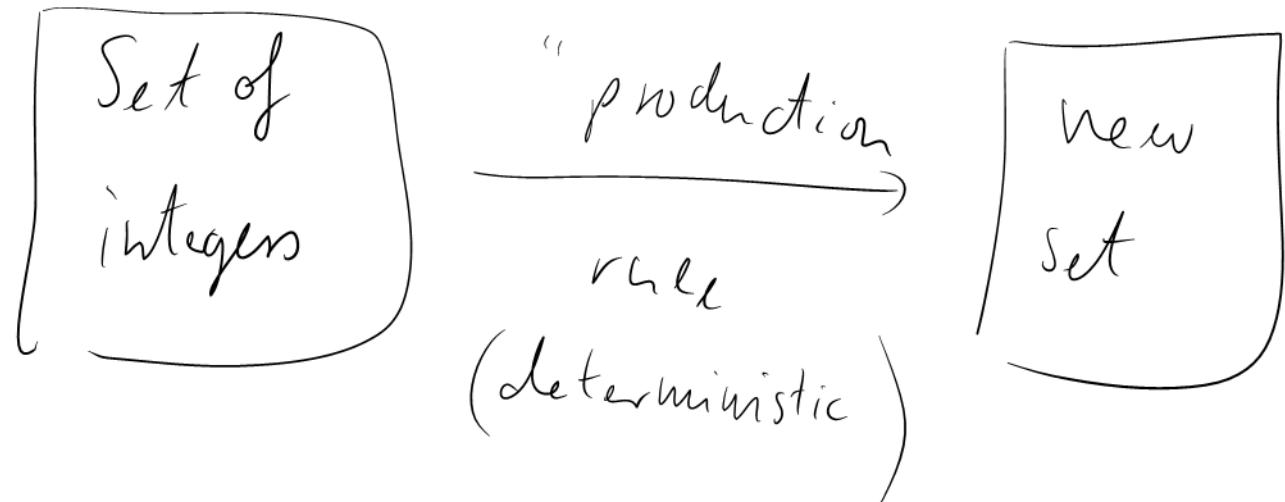


# Random Number Generators

Pseudo -



Lif. - Numerical Recipes

- Knuth, The Art of Computer Programming  
(vol. 2)

- Toms, Math. Comput., 19, 201 (1965)

→ Hner / BD / Fuchsberg, Comput. Phys. Comm. 103, 1 (1997)

✓ Banke / Mertus, PRE 75, 066701 (2007)

Simplest: "linear congruential"

$$x_{n+1} = (a \cdot x_n + c) \bmod m$$

popular choice:  $a = 7^5 = 16807$ ,  $c = 0$

$m = 2^{31} - 1$  (largest integer  
on 32-bit  
prime #)

this has maximum period

$$m-1 = 2^{31}-2 \approx \underline{2 \cdot 10^9}$$

rule of thumb: # of produced  $\approx N$

$\ll \sqrt{\text{period}}$

3 yrs CPU time  $\approx 10^8$  sec.

$10^6$  processes, one RN per machine cycle:

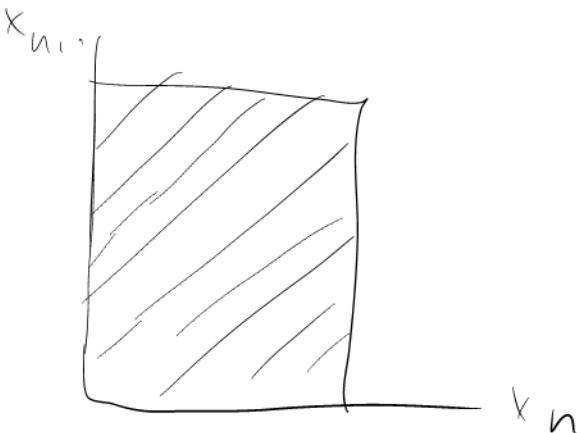
$$\# \text{ of RN} \approx 10^8 \cdot 10^9 \cdot 10^6 = \underline{\underline{10^{23}}}$$

period  $\gg 10^{46}$

correlations

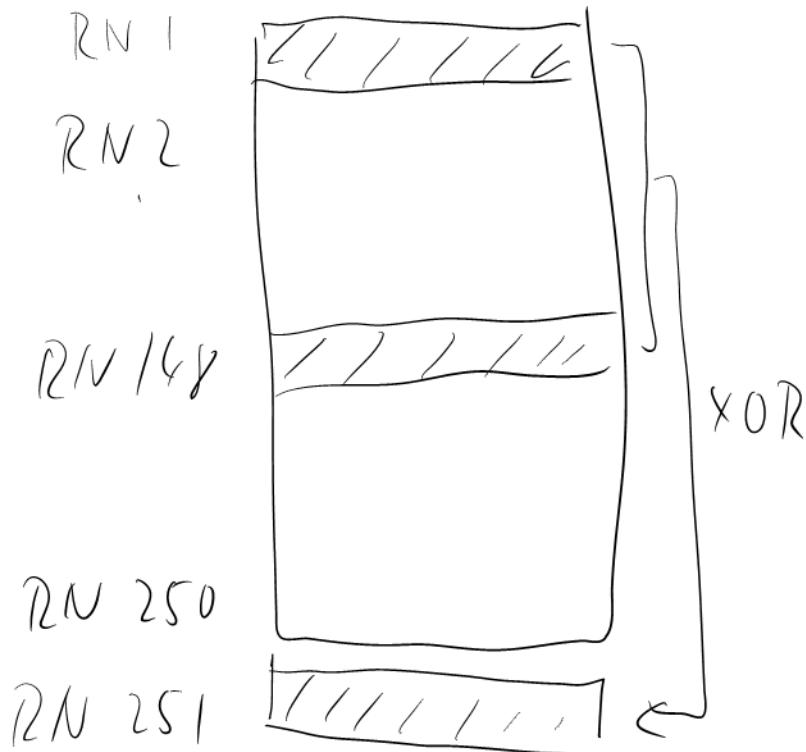
"d-space non-uniformity"

$$d=2$$



work on arrays of RNS

R250:



"shift-register RNS"  
triplet correlations

XOR: bitwise addition without carry

$$\emptyset \text{ XOR } \emptyset = \emptyset$$

$$1 \text{ XOR } 1 = 0$$

$$1 \text{ XOR } 0 = 1$$

$$0 \text{ XOR } 1 = 1$$

$$x_n = x_{n-10} \} \text{ XOR } x_{n-250}$$

altvn.:

$$x_n = x_{n-168} \text{ XOR } \underline{\underline{x_{n-521}}}$$

period :  $\equiv$  largest possible  $\cong 2^{250} \approx 2 \cdot 10^{75}$   
 $\Rightarrow 10^{46}$

combine LCG & S21 via XOR of the output

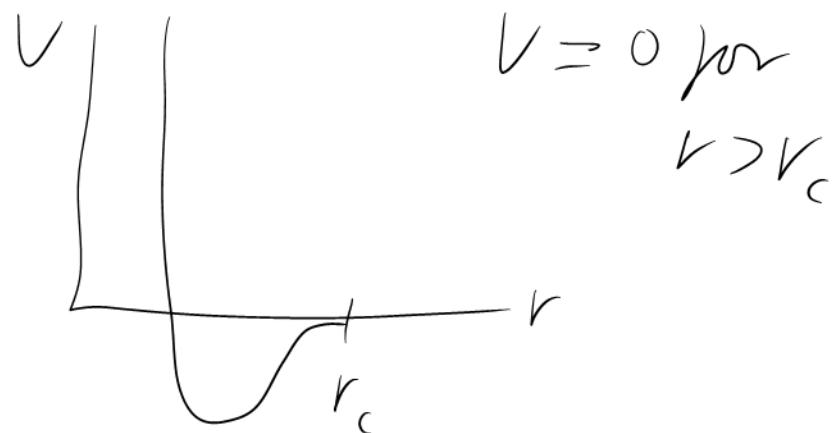
# Molecular Dynamics

$$m_i \ddot{\vec{r}}_i = \vec{F}_i = -\frac{\partial U}{\partial \vec{r}_i}$$

$$V_{LS} = \begin{cases} 4\epsilon \left[ \left(\frac{r}{r_s}\right)^{12} - \left(\frac{r}{r_s}\right)^6 \right] & r < 2.5r_s \\ 0 & \text{else} \end{cases}$$

typically:

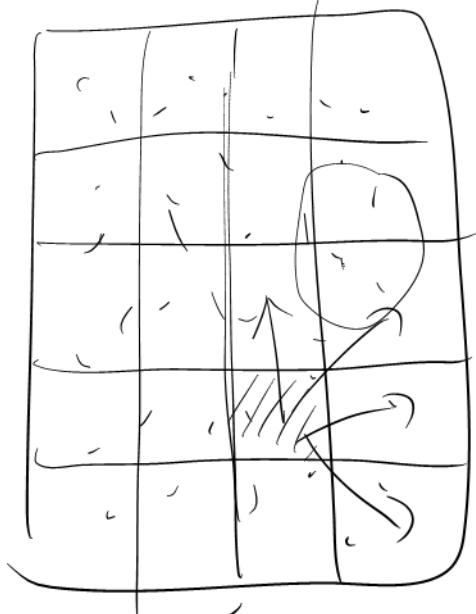
$$U = \sum_{i < j} V(|\vec{r}_i - \vec{r}_j|)$$



(i) how to calculate  $\vec{F}_i$ ?

- all pairs:  $O(N^2)$   $\frac{N(N-1)}{2}$

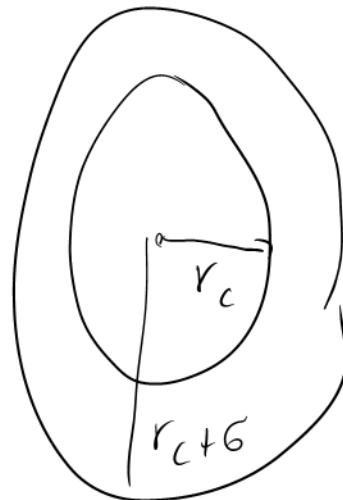
Short-range interactions



boxes:  $O(N)$  calculation

$r_c + \delta \rightarrow$  "skin"  $\delta > 0$

Verlet table:



Store all the neighbors  
within  $r_c + \delta$ , for each particle

$$\max_i \Delta r_i^2 > \frac{\delta^2}{4} : \text{UPDATE}$$

$\overset{\curvearrowleft}{G/2}$        $\overset{\curvearrowleft}{G/2}$   
before  
 $\curvearrowleft R_c + G$

after:

$\curvearrowleft R_c$

problem (ii)

Integrate Newton II!

useful: so-called "symplectic integrators"

based on Hamiltonian formalism

[Leimkuhler / Reich, Simulating Hamiltonian Dynamics]

$$\dot{q}_i = \frac{\partial \mathcal{H}}{\partial p_i} \quad \dot{p}_i = -\frac{\partial \mathcal{H}}{\partial q_i}$$

Liouville  
operator  
 $\downarrow$

$$\frac{d}{dt} A = \sum_i \left( \underbrace{\frac{\partial \mathcal{H}}{\partial p_i} \frac{\partial}{\partial q_i} - \frac{\partial \mathcal{H}}{\partial q_i} \frac{\partial}{\partial p_i}}_{i \in q} \right) A = i \mathcal{L} A$$

$$i\ell = i\ell_q + i\ell_p \quad \underbrace{\text{Specialization:}}$$

$$\mathcal{H} = T(\{p_i\}) + U(\{q_i\})$$

$$\underline{i\ell_q} = \sum_i \frac{\partial T}{\partial p_i} \frac{\partial}{\partial q_i}$$

$\sum \uparrow$

$$i\ell_p = \sum_i \left( -\frac{\partial U}{\partial q_i} \right) \frac{\partial}{\partial p_i}$$

$\uparrow \sum \uparrow$

$$= \sum_i \frac{\partial}{\partial q_i} \frac{\partial T}{\partial p_i}$$

$$= - \sum_i \frac{\partial}{\partial p_i} \frac{\partial U}{\partial q_i}$$

$\ell$  is self-adjoint, but also  $\ell_q$  and  $\ell_p$

separately

$$\begin{aligned} (A | \ell_q B) &= \frac{1}{i} (A | ; \ell_q B) = \frac{1}{i} \int d\Gamma A^*(\Gamma) \sum_i \frac{\partial T}{\partial p_i} \frac{\partial}{\partial q_i} B(\Gamma) \\ &= -\frac{1}{i} \int d\Gamma B(\Gamma) \sum_i \frac{\partial T}{\partial p_i} \frac{\partial}{\partial q_i} A^*(\Gamma) = -\frac{1}{i} \int d\Gamma B(\Gamma) ; \ell_q A^*(\Gamma) \\ &= \underbrace{\int d\Gamma B(\Gamma)}_{\ell_q^*} (-\ell_q) A^*(\Gamma) = \int d\Gamma B(\Gamma) [\ell_q A]^* = \\ &\quad \ell_q^* = (\ell_q A | B) \\ \ell_p : \text{ analogous,} \end{aligned}$$

$e^{i\ell^+}$  unitary,  $e^{i\ell^-}$  unitary,  $e^{i\ell^P}$  unitary

would like to know:

$$e^{i\ell^+ q_i}, e^{i\ell^+ p_i} = ? \quad \text{HARD}$$

$$e^{i\ell^- q_i} = ? \quad e^{i\ell^P q_i} = ?$$

$$e^{i\ell^- p_i} = ? \quad e^{i\ell^P p_i} = ? \quad \underline{\text{EASY}}$$

$$\exp(i\ell_q t) = \sum_{n=0}^{\infty} \frac{t^n}{n!} (i\ell_q)^n = \sum_{n=0}^{\infty} \frac{t^n}{n!} \left\{ \sum_j \frac{\partial T}{\partial p_j} \frac{\partial}{\partial q_j} \right\}^n$$

$$\exp(i\ell_q t) q_i = q_i + t \frac{\partial T}{\partial p_i} = q_i + v_i t$$

$$\exp(i\ell_q t) p_i = p_i$$

$$\exp(i\ell_p t) q_i = q_i$$

$$\exp(i\ell_p t) p_i = p_i - \frac{\partial U}{\partial q_i} t = p_i + F_i t$$

operator-splitting

time step  $h$ ,  $h \rightarrow 0$

$$\exp[i(\ell_q + \ell_p)h] \simeq \exp\left\{i\ell_p \frac{h}{2}\right\} \exp\left\{i\ell_q h\right\} \exp\left\{i\ell_p \frac{h}{2}\right\}$$

$\ell$  + error:  $O(h^3)$  or smaller

Taylor wrt.  $L$ :

agreement up to  $O(h^2)$  inclusively

- conserves the phase-space volume
- $h \rightarrow -h$ : time-reversal symmetric

## algorithm:

$$p_i\left(\frac{h}{2}\right) = p_i(0) + \bar{f}_i(0) \frac{h}{2}$$

$$q_i(h) = q_i(0) + v_i\left(\frac{h}{2}\right) h$$

$$p_i(h) = p_i\left(\frac{h}{2}\right) + \bar{f}_i(h) \frac{h}{2}$$

velocity

Verlet

time step :  $10^{-2} \times$  shortest oscillation time

$\log \delta_1^2$



ergodic  $t^2$

cage

$\log \epsilon$

$t^2$   $t^2$